

# A Tutte-type canonical decomposition of 3- and 4-connected graphs

Jan Kurkofka

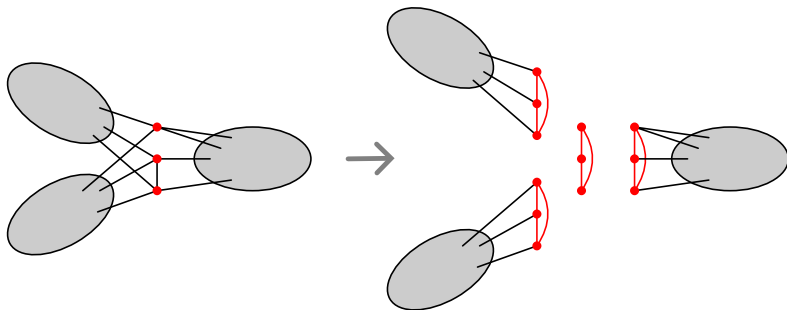


Joint work with Tim Planken

Ukrainian Summer School in Combinatorics 2025

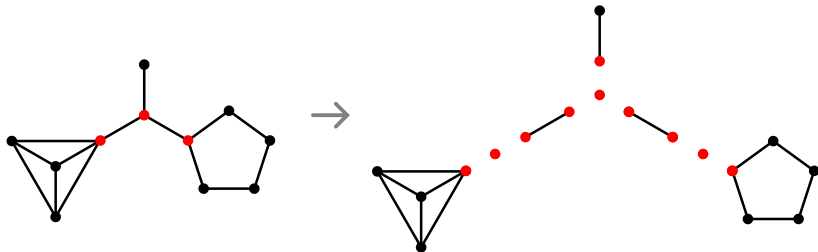
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into parts that are  $(k + 1)$ -con'd or 'basic'.

Decomposing  $G$  along a  $k$ -separator:



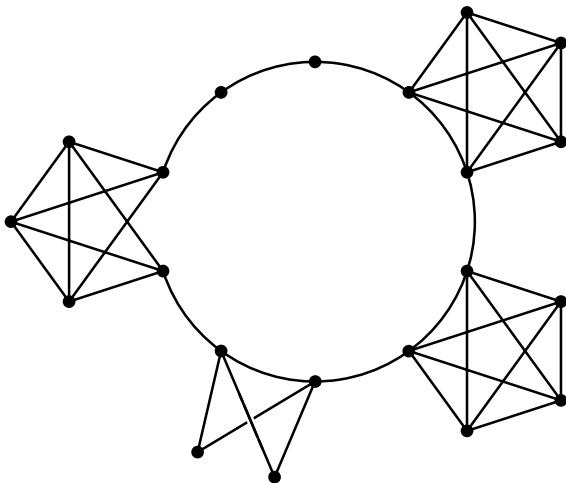
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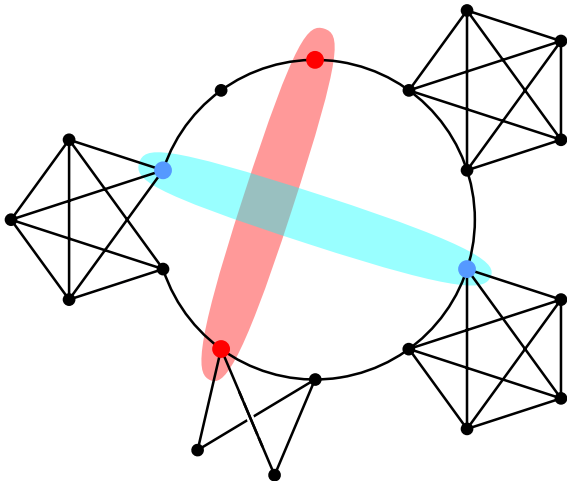
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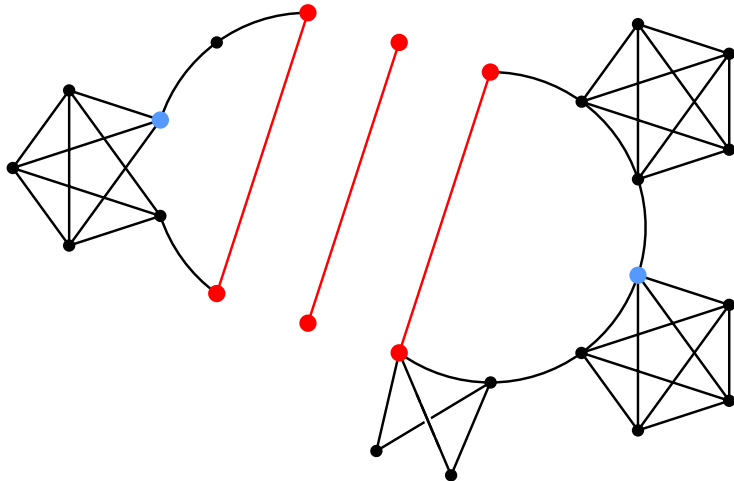
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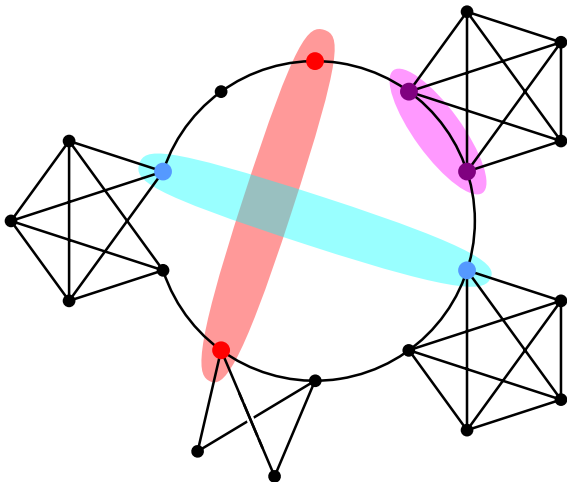


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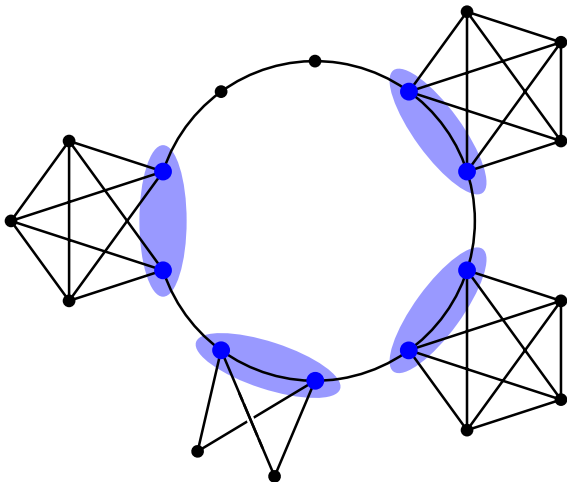


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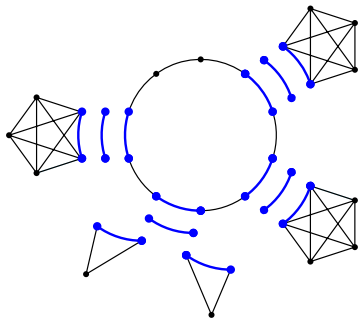
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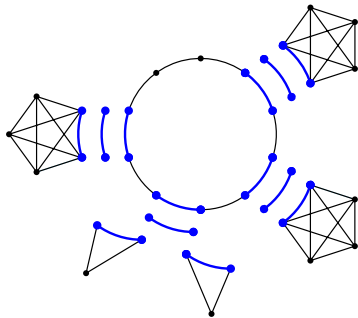
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Theorem (Tutte 66), SPQR-trees

Every 2-con'd  $G$  decomposes along its totally-nested 2-separators  
into 3-con'd graphs, cycles and  $K_2$ 's.

- **canonical**: isomorphisms map parts to parts
- **tree-decomposition** (for fans)

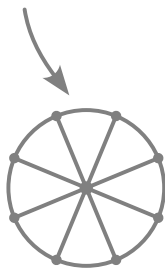


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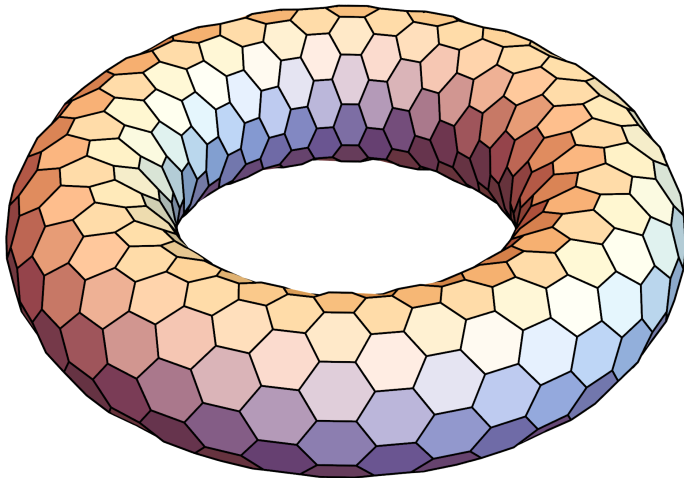


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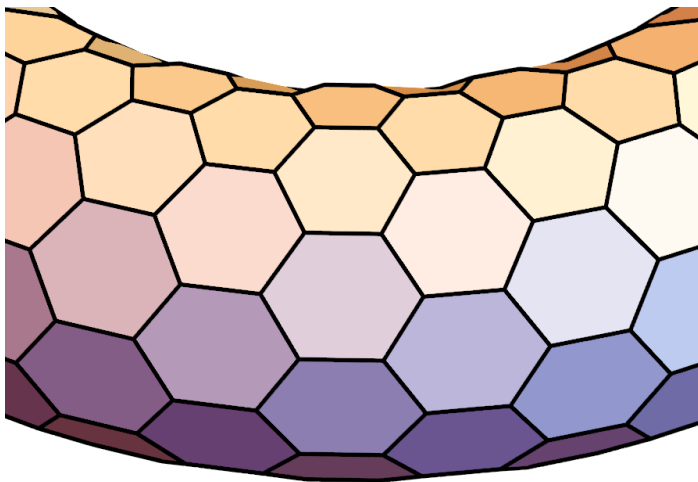
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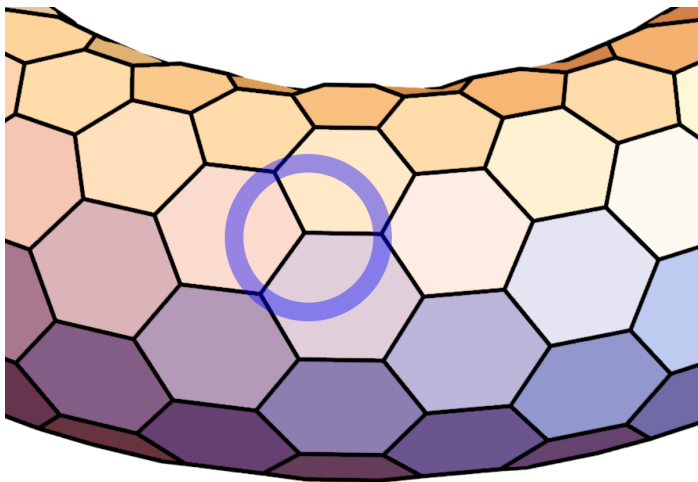
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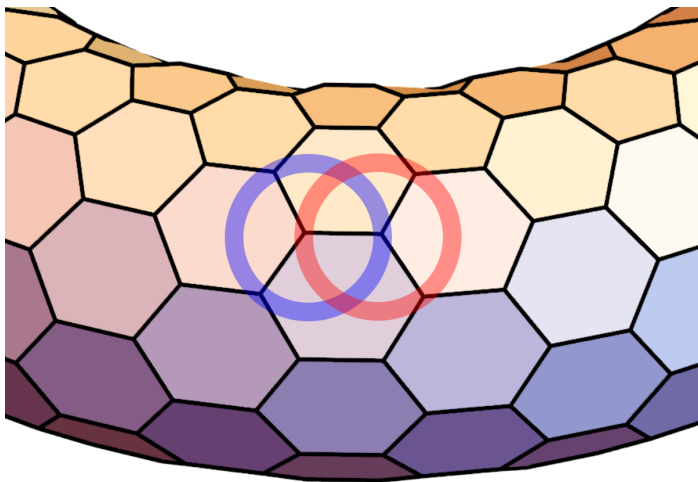
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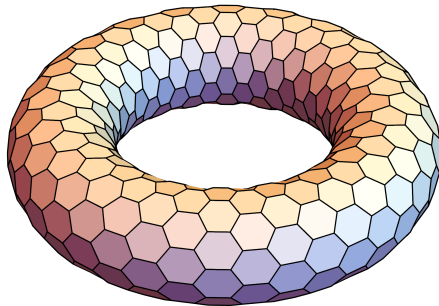
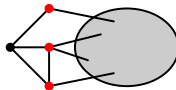


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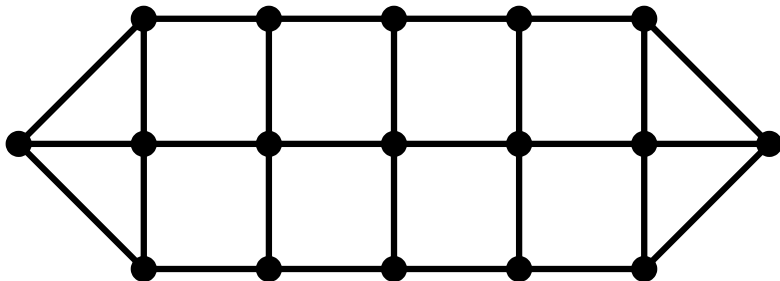


3-con'd and every 3-separator has form



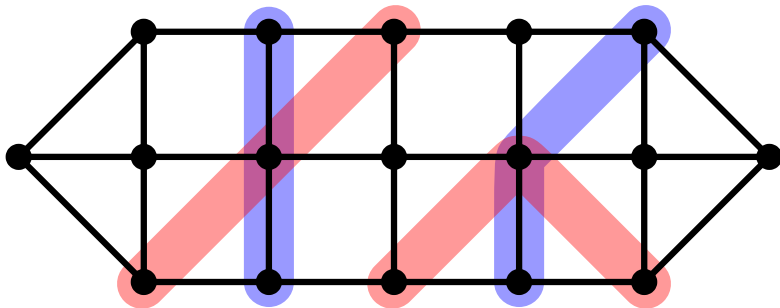
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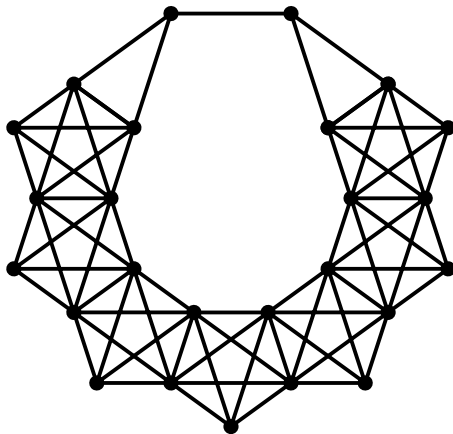
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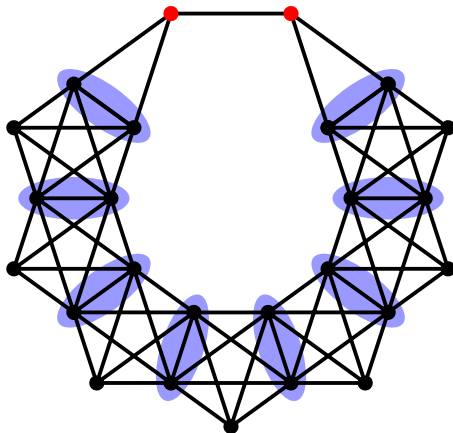
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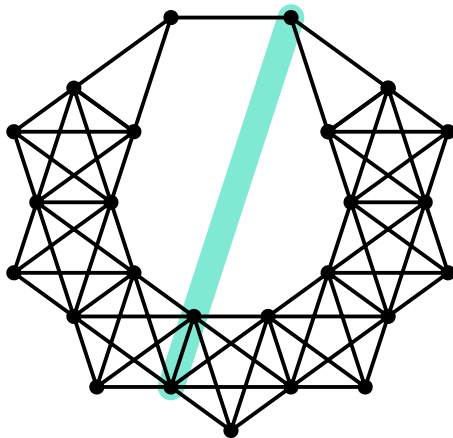
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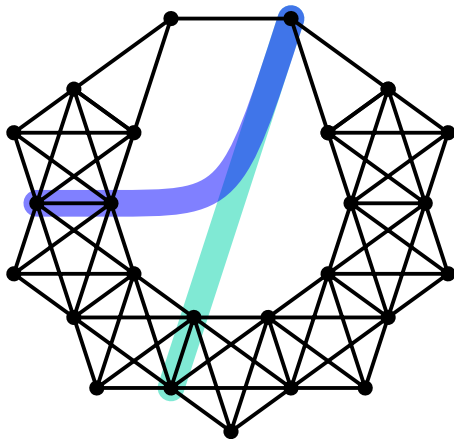
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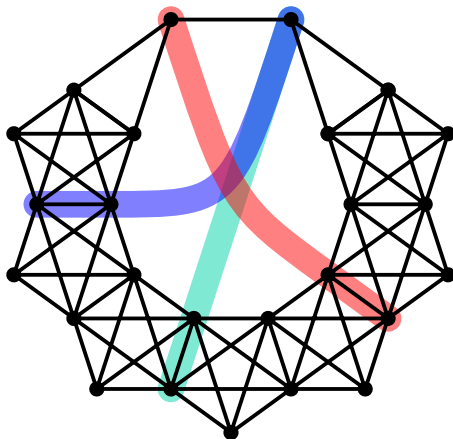
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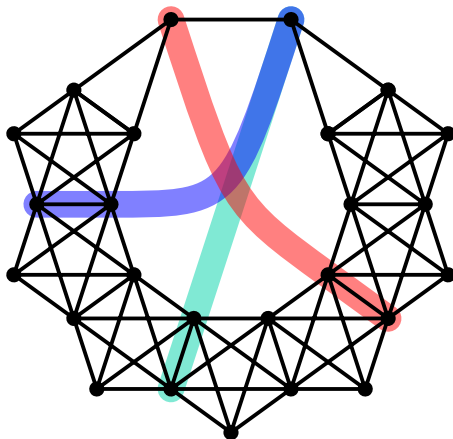
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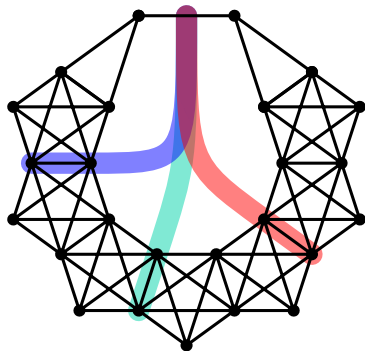
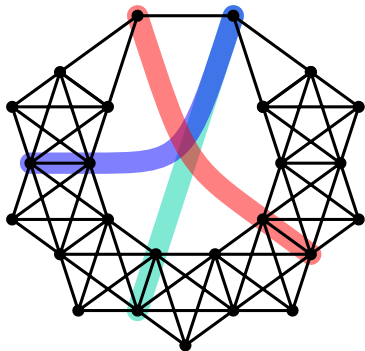
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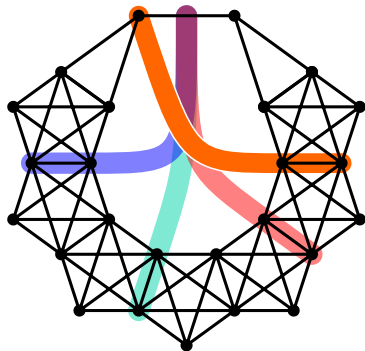
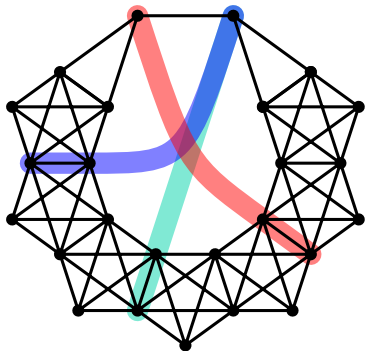
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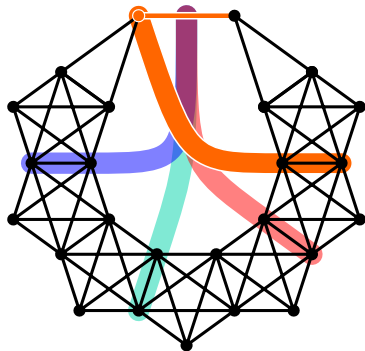
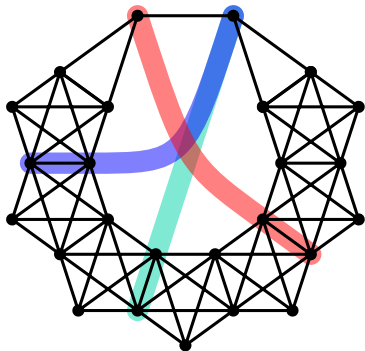
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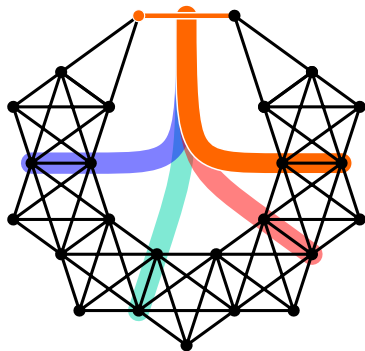
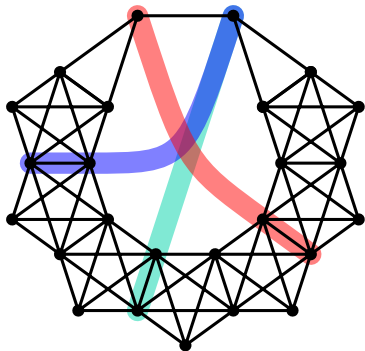
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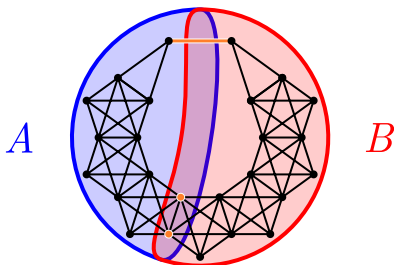
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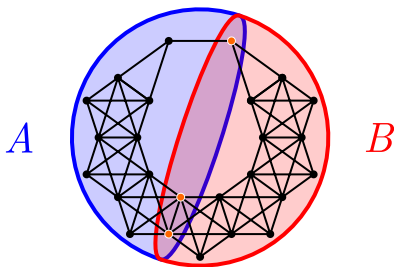




*mixed-separation* of  $G$ :  $(A, B)$  with  $A \cup B = V(G)$  and  $A, B \neq V(G)$

*separator* of  $(A, B)$ :  $(A \cap B) \cup E(A \setminus B, B \setminus A)$

A **tri-separation** of  $G$  is a mixed-sep'n  $(A, B)$  with  $|\text{sep}'r| = 3$  s.t. every  $v_x$  in  $A \cap B$  has  $\geq 2$  neighb's in  $A$  and  $B$ .

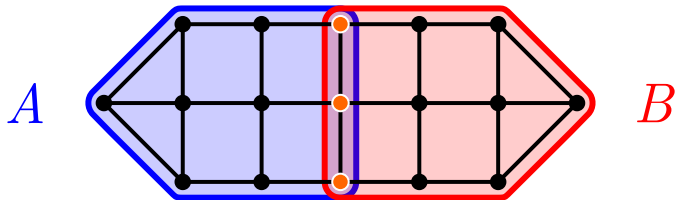


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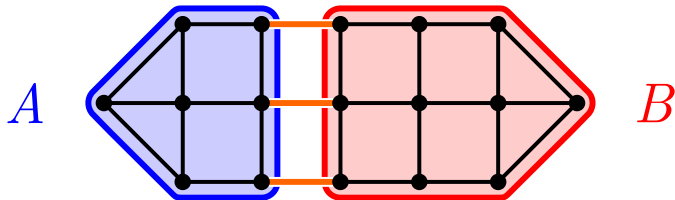




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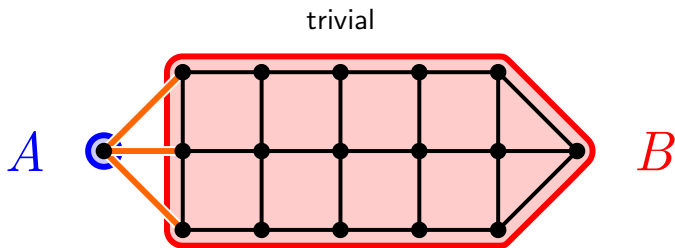
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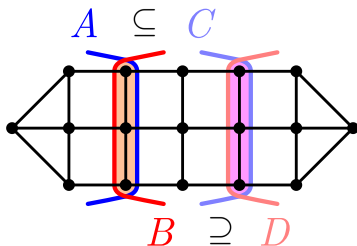
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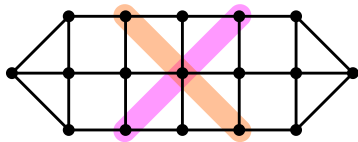
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$(A, B)$  and  $(C, D)$  are *nested* if  $A \subseteq C$  and  $B \supseteq D$   
 after possibly switching  $A$  with  $B$  or  $C$  with  $D$ ;  
 otherwise they *cross*.

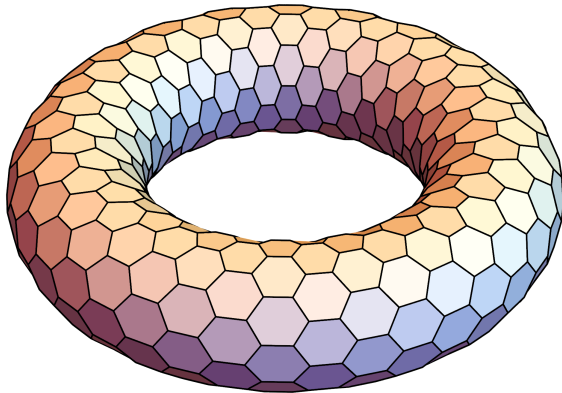
nested



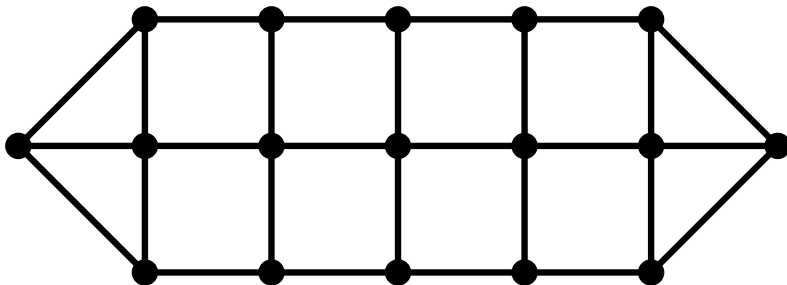
crossing



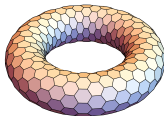
totally-nested nontrivial tri-separations



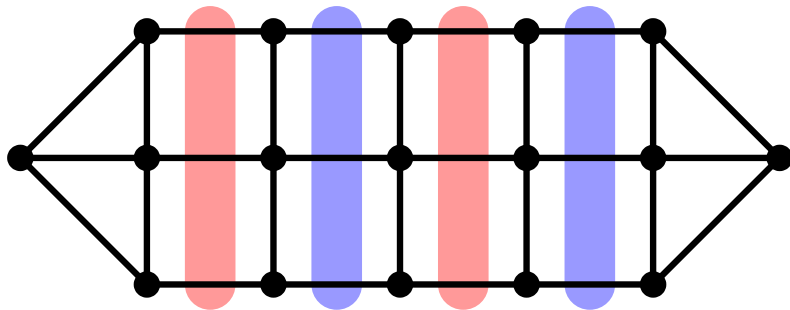
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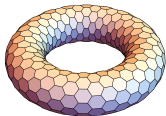
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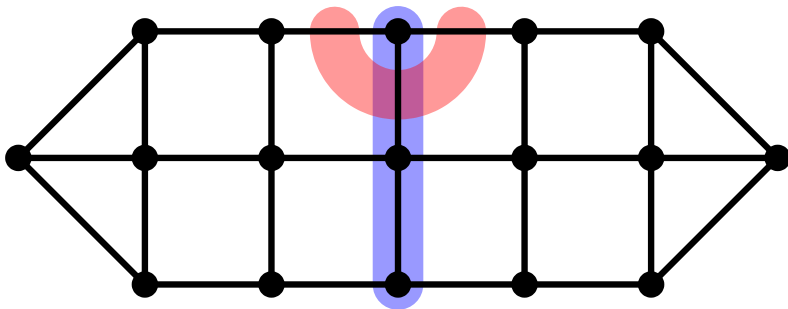
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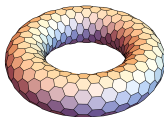
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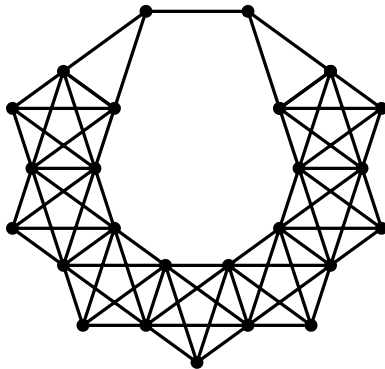


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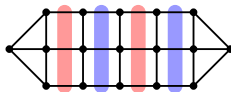
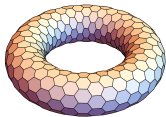




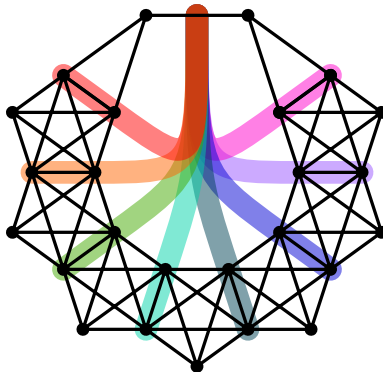
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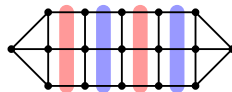
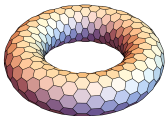
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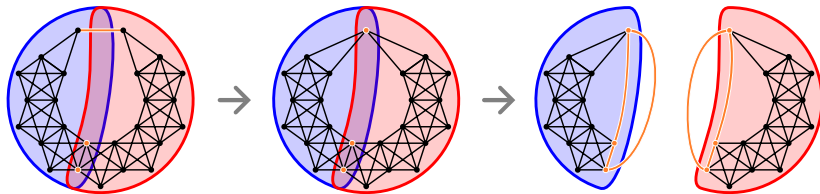
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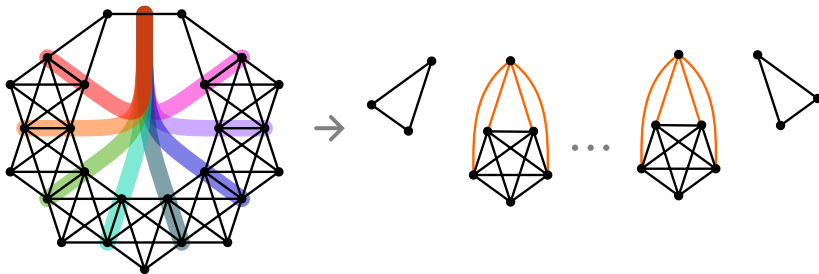
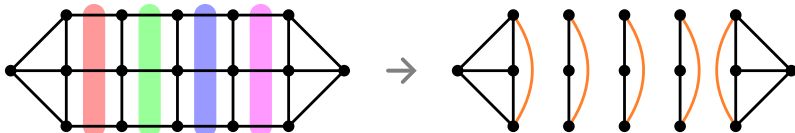
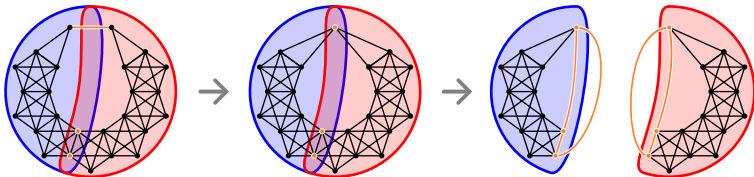


none



Decomposing along a tri-separation

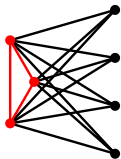




Theorem (Carmesin & K. 23)

Every 3-con'd  $G$  decomposes along its totally-nested nontrivial tri-separations into parts that are

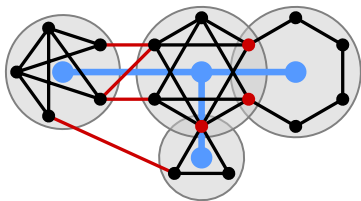
- quasi 4-con'd
- wheels
- thickened  $K_{3,m}$



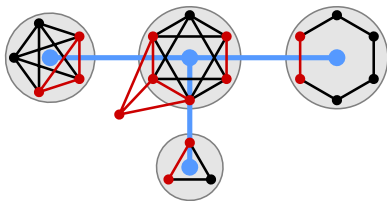
$3$        $m$

- canonical ✓
- for tree-decomposition fans:

mixed-tree-decomposition

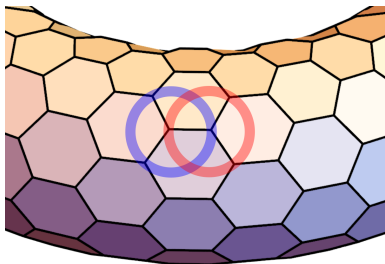


torsos



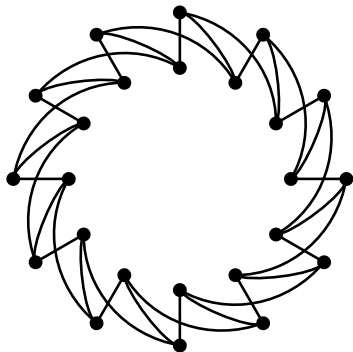
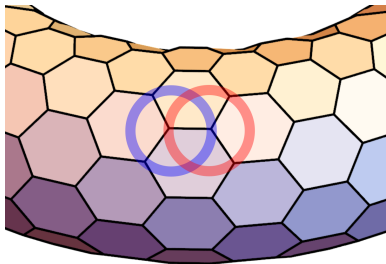
$$k \geq 4?$$

Challenge 1: Torus-example for  $k \geq 3$

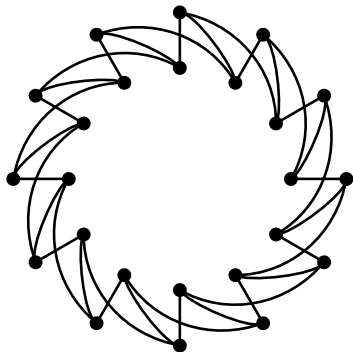
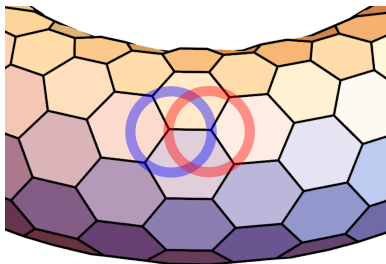




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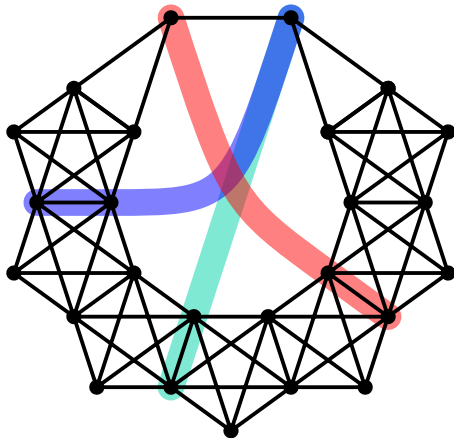


Challenge 1: Torus-example for  $k \geq 3$

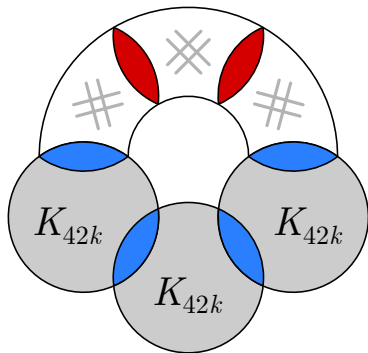


quasi- $k$ -connected:  $k$ -con'd and every  $k$ -sep'r cuts off  $\leq 1$  vertex

## Challenge 2



## Challenge 2

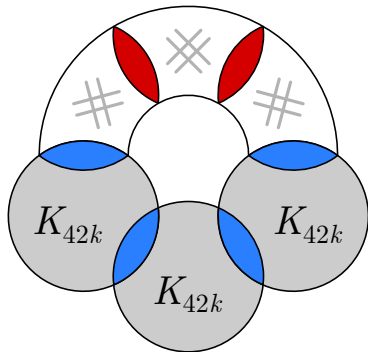


$k$  odd

red:  $\lfloor k/2 \rfloor$ -cliques

blue:  $\lceil k/2 \rceil$ -cliques

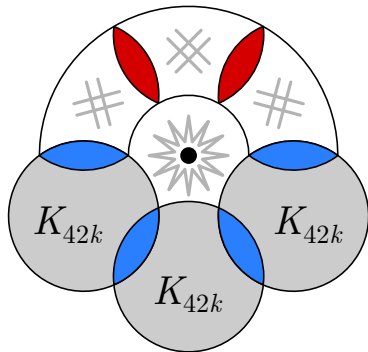
## Challenge 2



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red:  $\lfloor k/2 \rfloor$ -cliques

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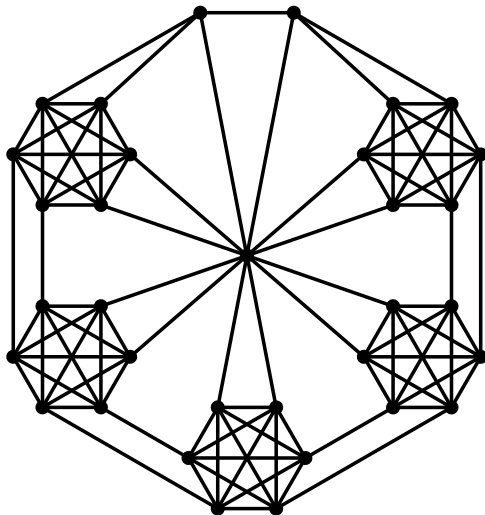


$k$  even

red:  $(\frac{k}{2} - 1)$ -cliques

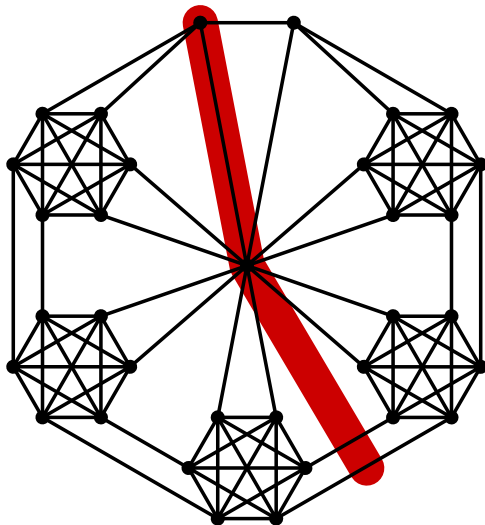
blue:  $\frac{k}{2}$ -cliques

## Challenge 2



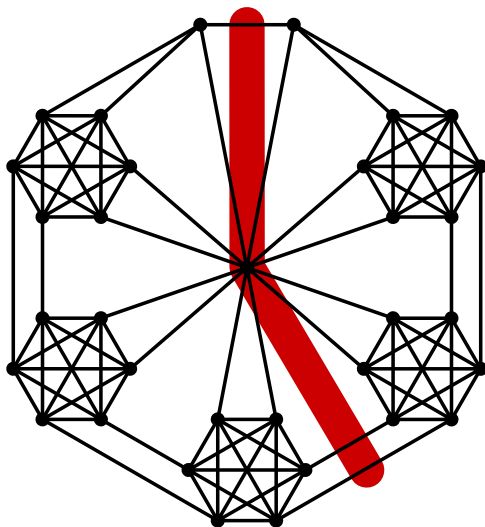
## Challenge 2

Verbatim extension of tri-sep'ns to  $k = 4$ :



## Challenge 2

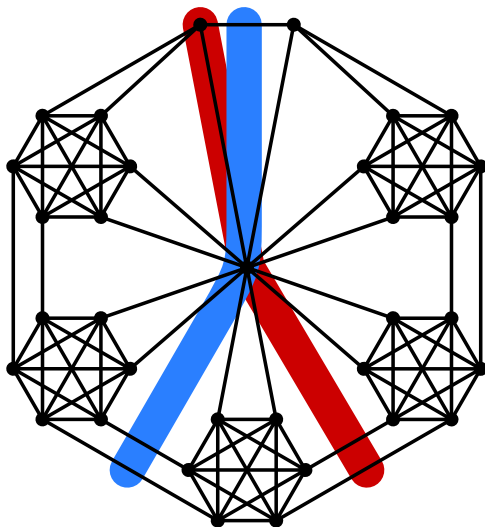
Verbatim extension of tri-sep'ns to  $k = 4$ :



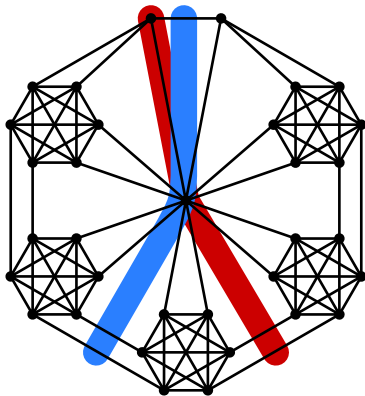


## Challenge 2

Verbatim extension of tri-sep'ns to  $k = 4$ :

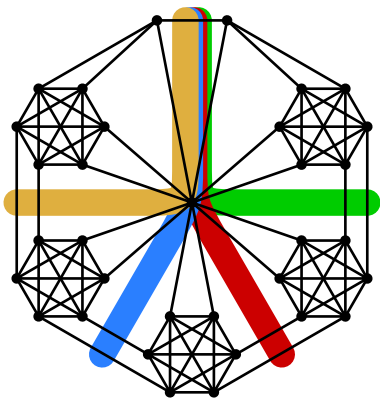






A **tetra-separation** is a mixed-sep'n  $(A, B)$  with  $|\text{sep}'r| = 4$  s.t.:

- every  $v_x$  in  $A \cap B$  has  $\geq 2$  neighb's in  $A \setminus B$  and  $B \setminus A$
- the edges in the sep'r form a matching



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Main result (K. & Planken 25)

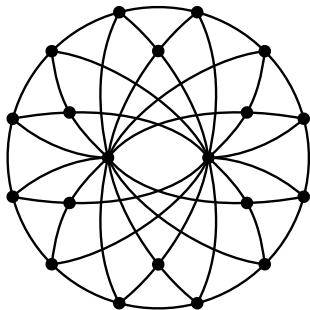
Every 4-con'd  $G$  decomposes along its totally-nested tetra-separations into parts that are

- quasi 5-con'd

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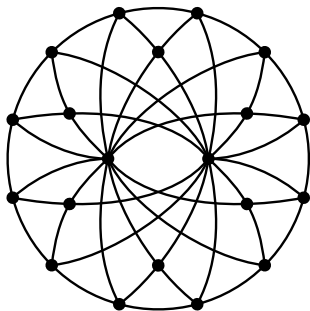
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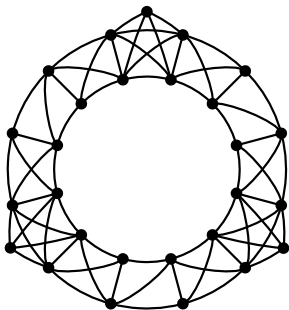
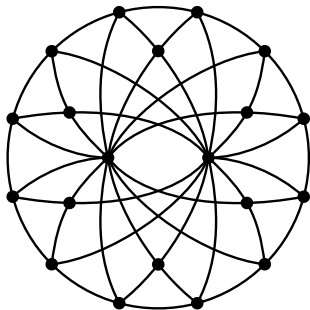
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Every 4-con'd  $G$  decomposes along its totally-nested tetra-separations into parts that are

- quasi 5-con'd
- thickened  $K_{4,m}$
- generalised double-wheels
- cycle of triangles and 3-con'd graphs on  $\leq 5$  vxs.





con'd

block-cut

con'd



block-cut

con'd



2-con'd

$K_2, K_1$

block-cut

con'd



2-con'd

$K_2, K_1$

Tutte



block-cut

con'd



2-con'd

$K_2, K_1$

Tutte



3-con'd

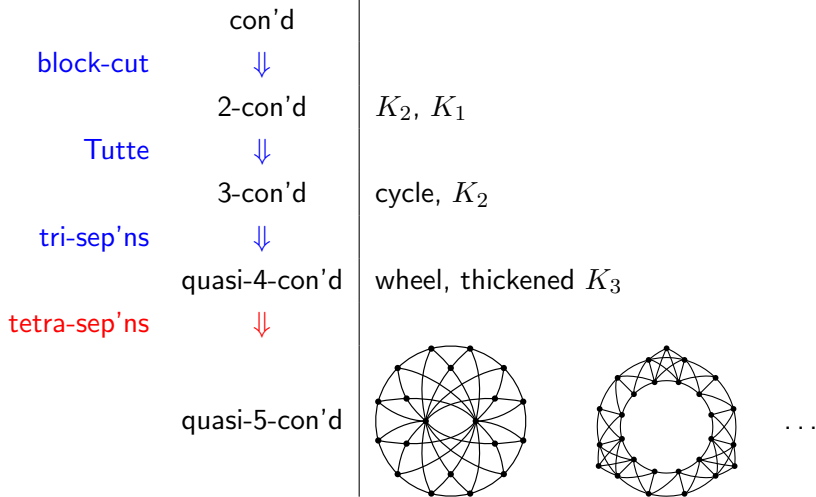
cycle,  $K_2$

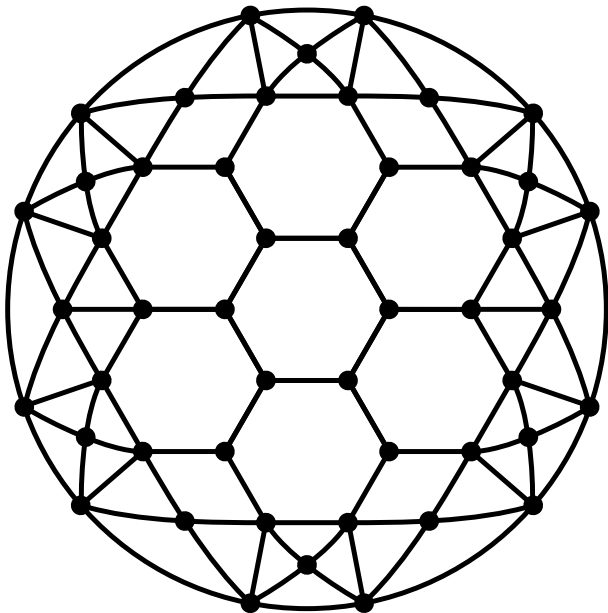
	con'd	
block-cut	$\Downarrow$	
	2-con'd	$K_2, K_1$
Tutte	$\Downarrow$	
	3-con'd	cycle, $K_2$
tri-sep'ns	$\Downarrow$	

	con'd	
block-cut	$\Downarrow$	
	2-con'd	$K_2, K_1$
Tutte	$\Downarrow$	
	3-con'd	cycle, $K_2$
tri-sep'ns	$\Downarrow$	
	quasi-4-con'd	wheel, thickened $K_3$

	con'd	
block-cut	⇓	
	2-con'd	$K_2, K_1$
Tutte	⇓	
	3-con'd	cycle, $K_2$
tri-sep'ns	⇓	
	quasi-4-con'd	wheel, thickened $K_3$
tetra-sep'ns	⇓	

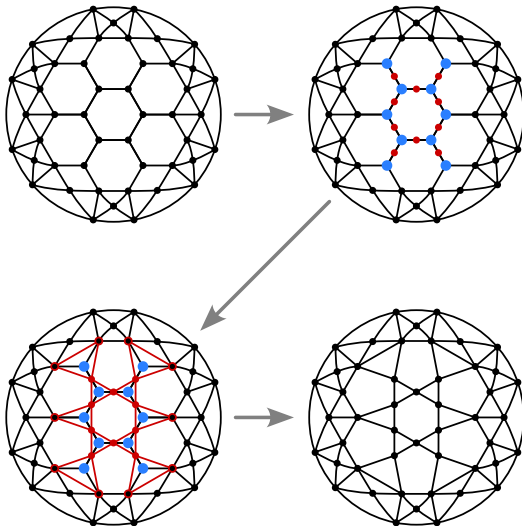




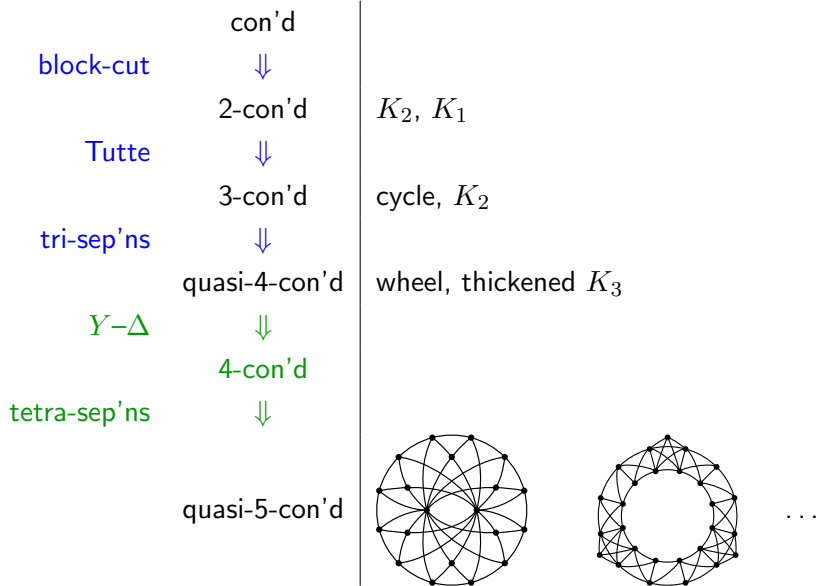


# canonical $Y-\Delta$ transformation

quasi-4-con'd



4-con'd



Problem: Classify all **vertex-transitive** finite con'd  $G$

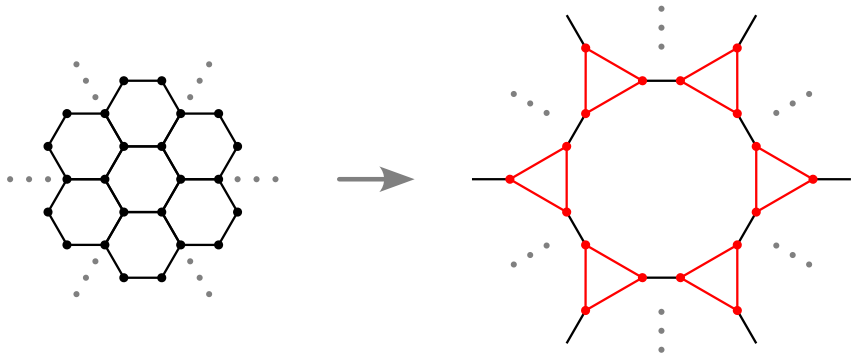
Problem: Classify all **vertex-transitive** finite con'd  $G$

Approach: Low connectivity first

Theorem (Carmesin & K. 23)

Every **vertex-transitive** finite con'd  $G$  is either

- a cycle,  $K_2$ ,  $K_1$ ,
- quasi-4-con'd, or
- $K_3$ -expansion of a quasi-4-con'd 3-regular arc-transitive graph.



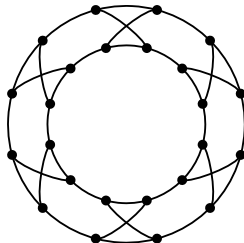
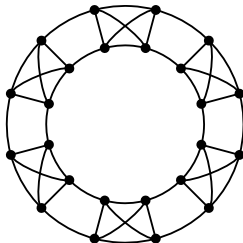
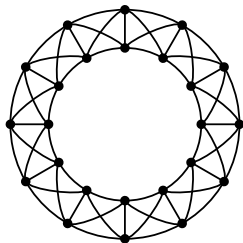
Theorem (K. & Planken 25)

Every **quasi-4-con'd vertex-transitive** finite  $G$  is either

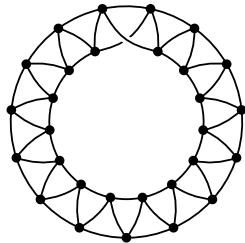
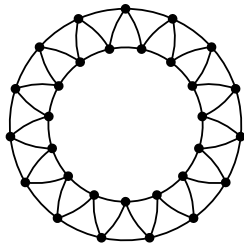
- bagel-like,
- cube-like,
- quasi-5-con'd, or
- $K_4/C_4$ -expansion of quasi-5-con'd 4-regular arc-transitive graph.



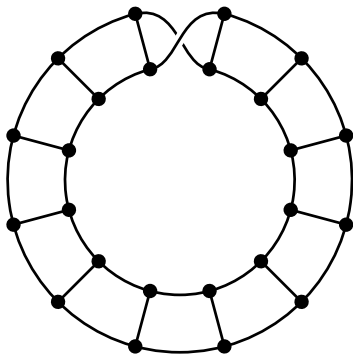
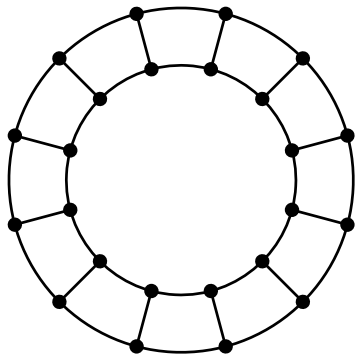
bagel-like



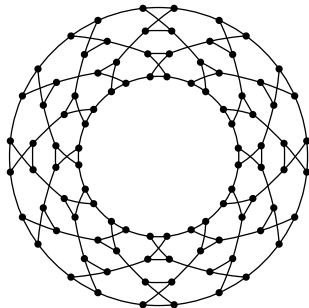
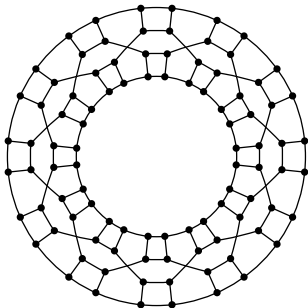
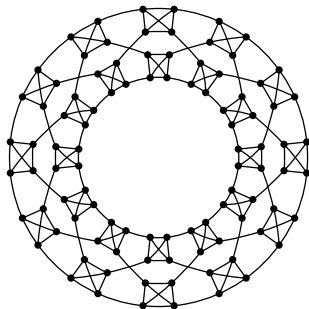
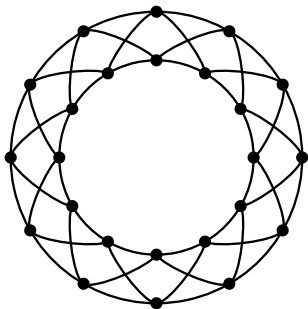
bagel-like



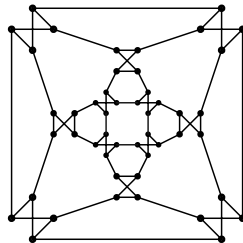
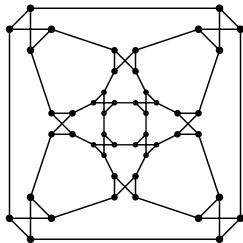
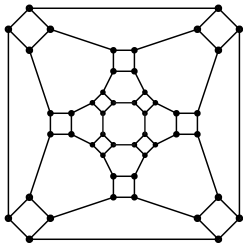
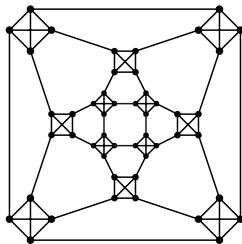
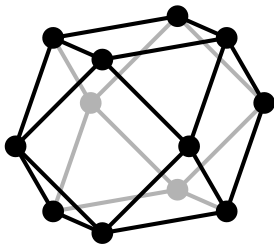
bagel-like



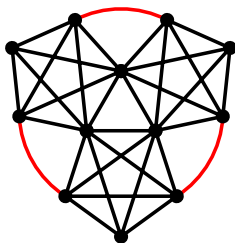
bagel-like



cube-like



Application: Connectivity Augmentation from 0 to 4



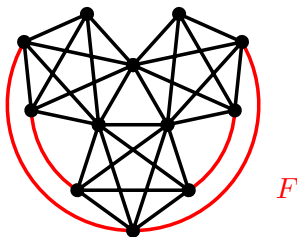
Theorem (Carmesin & Sridharan 25+)

$\exists$  FPT-algorithm with runtime  $C(\ell) \cdot \text{Poly}(|V(G)|)$  and

Input: Graph  $G$ ,  $\ell \in \mathbb{N}$  and  $F \subseteq E(\overline{G})$

Output: No, or  $\leq \ell$ -sized  $X \subseteq F$  such that  $G + X$  is 4-con'd

Application: Connectivity Augmentation from 0 to 4



Theorem (Carmesin & Sridharan 25+)

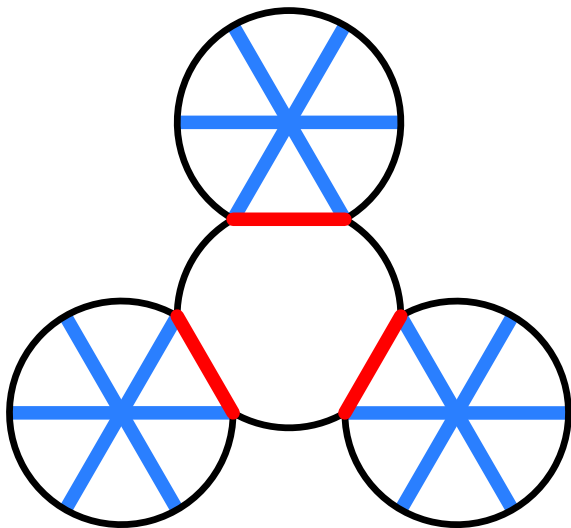
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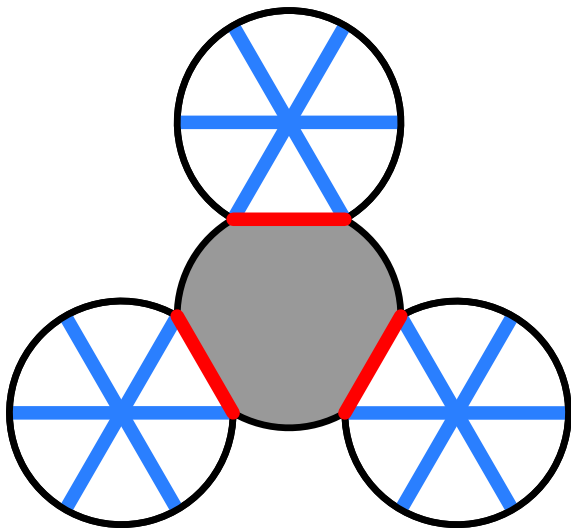
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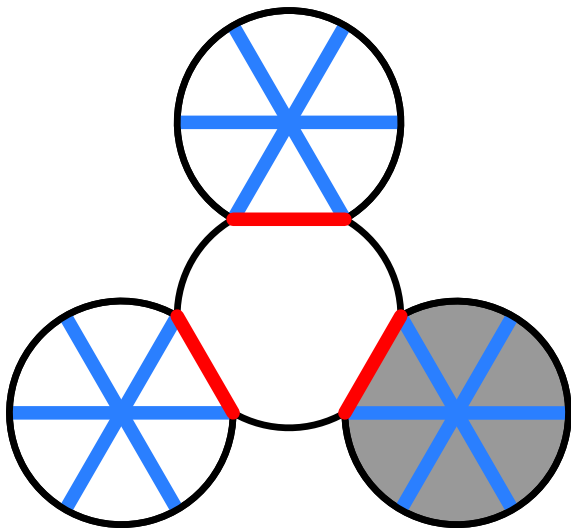
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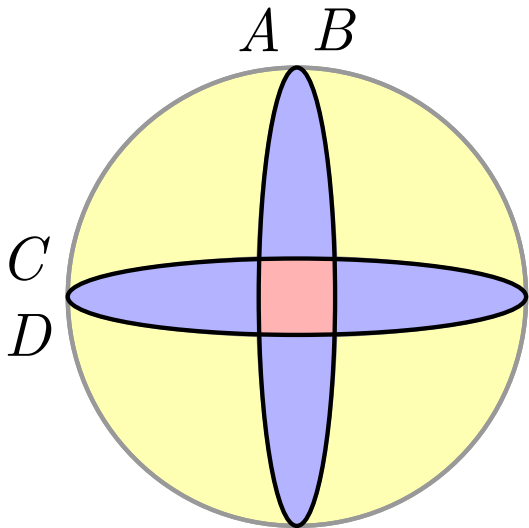
## Proof of tetra-decomposition

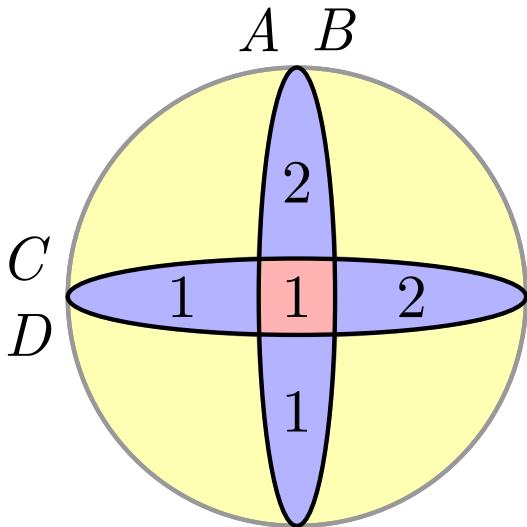




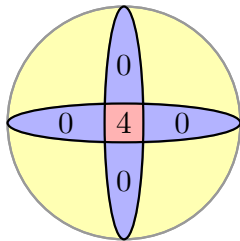
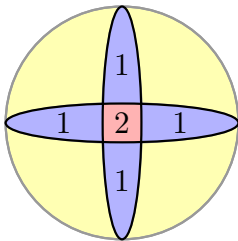
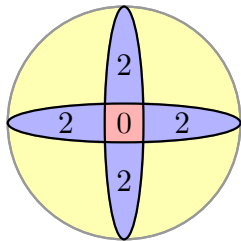




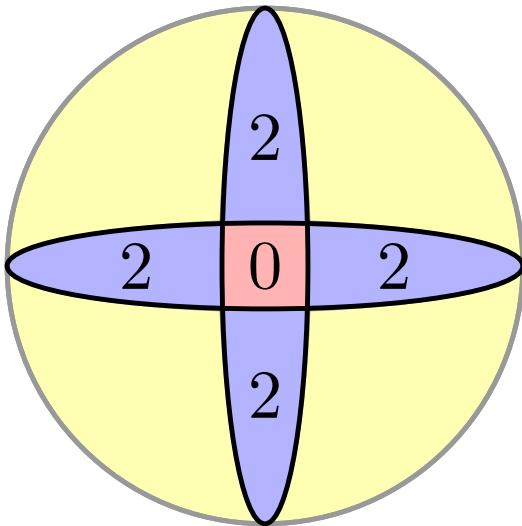


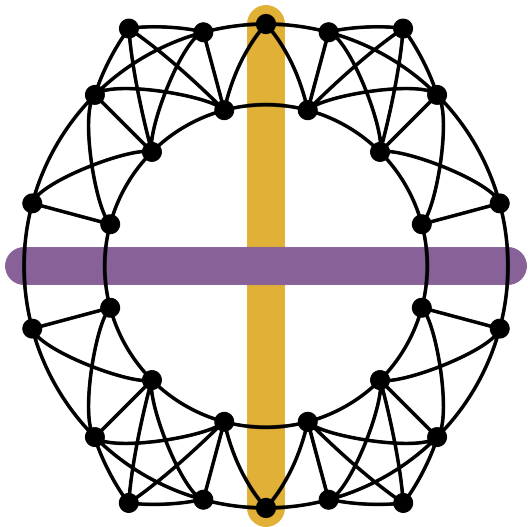


Crossing Lemma. Tetra-sep'ns only cross symmetrically:

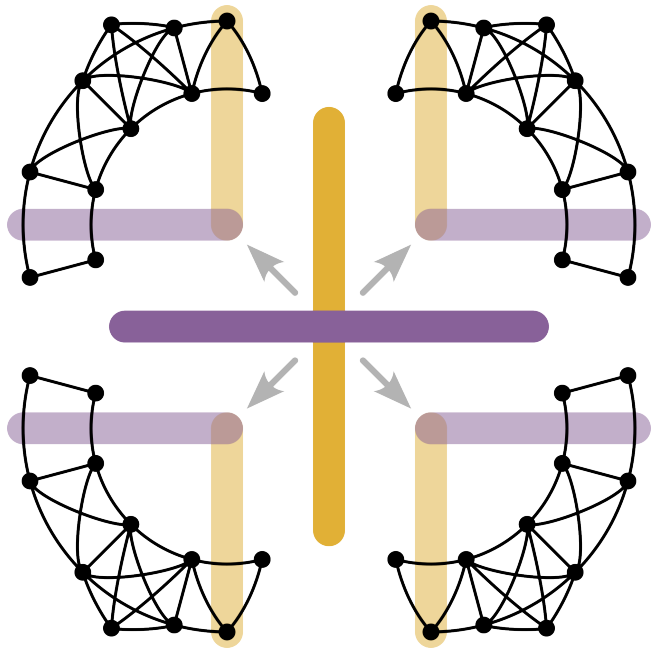


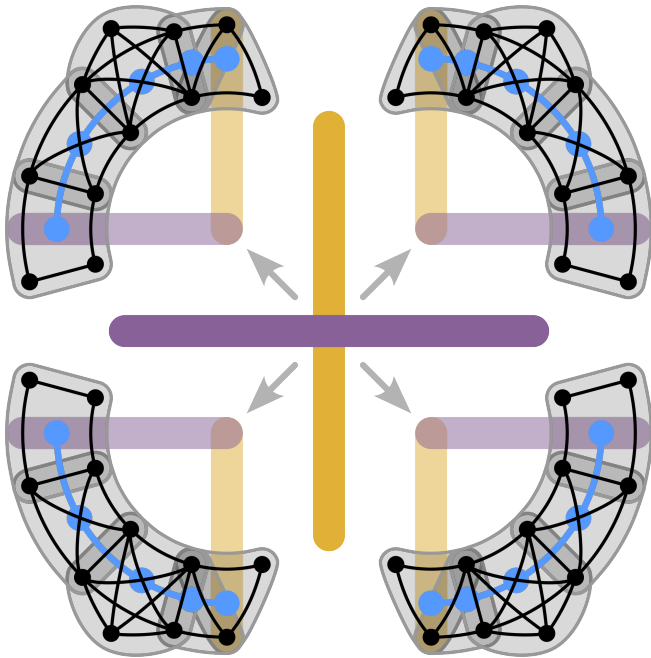
focus

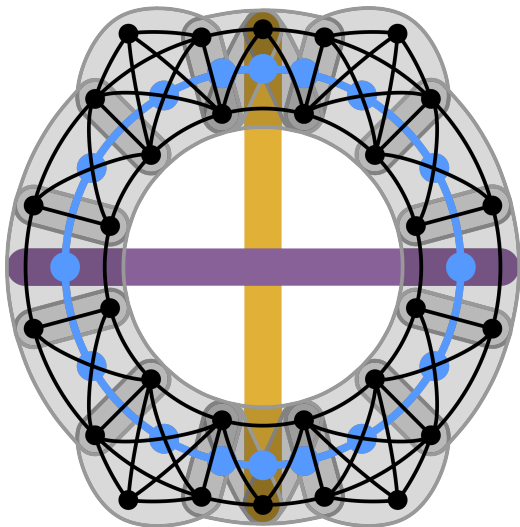


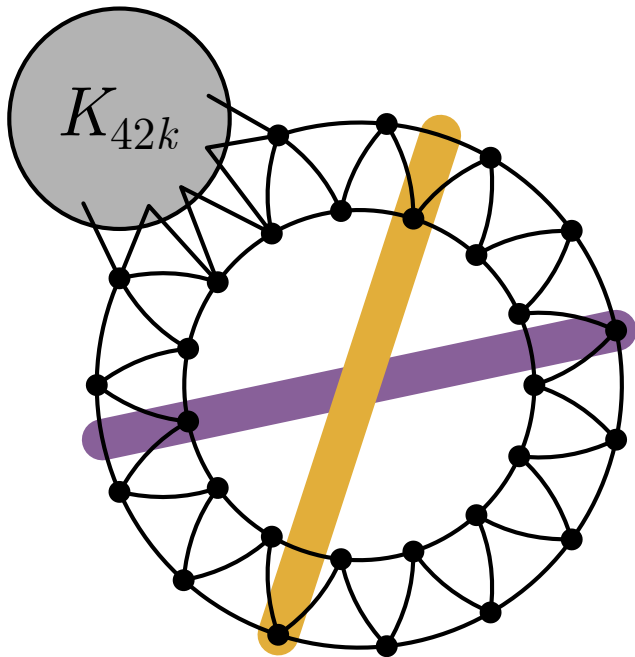


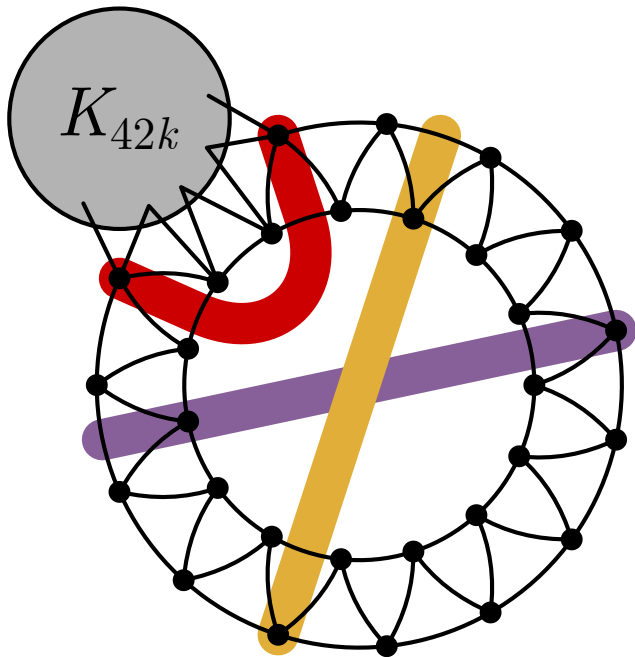






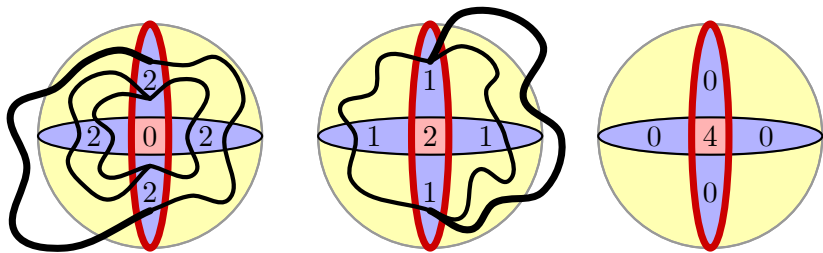


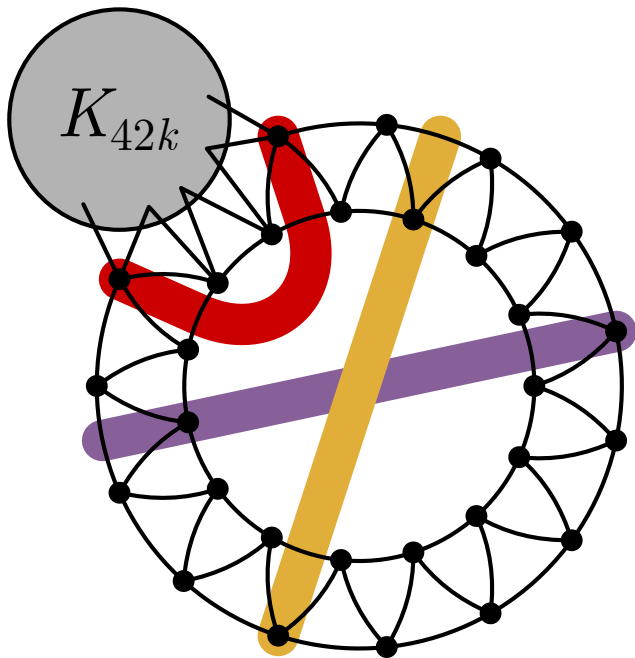


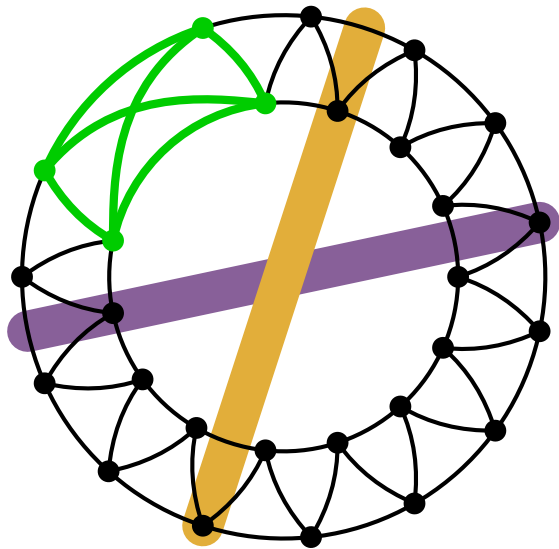


$(A, B)$  totally-nested

$\iff$  the sep'r of  $(A, B)$  is highly con'd:









Open: Extend the main result to *all*  $k$ .

Open: Directed graphs?

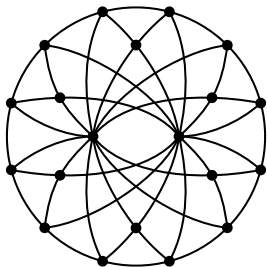
$k = 1$ : Bowler, Gut, Hatzel, Kawarabayashi, Muzi, Reich 23

$k \geq 2$ : ???

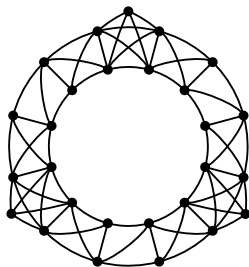
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Main result (K. & Planken 25)

Every 4-con'd  $G$  decomposes along its totally-nested tetra-sep'ns into parts that are quasi-5-con'd, thickened  $K_{4,m}$ 's,



or



Open: Graphs for  $k \geq 5$ . Digraphs for  $k \geq 2$ .